

# Chapter 7

## Concentration of Random Variables – Chernoff’s Inequality

By Sarel Har-Peled, April 12, 2018<sup>①</sup>

### 7.1. Concentration of mass and Chernoff’s inequality

#### 7.1.1. Example: Binomial distribution

Consider the binomial distribution  $\text{Bin}(n, 1/2)$  for various values of  $n$  as depicted in Figure 7.1 – here we think about the value of the variable as the number of heads in flipping a fair coin  $n$  times. Clearly, as the value of  $n$  increases the probability of getting a number of heads that is significantly smaller or larger than  $n/2$  is tiny. Here we are interested in quantifying exactly how far can we divert from this expected value. Specifically, if  $X \sim \text{Bin}(n, 1/2)$ , then we would be interested in bounding the probability  $\mathbb{P}[X > n/2 + \Delta]$ , where  $\Delta = t\sigma_X = t\sqrt{n}/2$  (i.e., we are  $t$  standard deviations away from the expectation). For  $t > 2$ , this probability is roughly  $2^{-t}$ , which is what we prove here.

More surprisingly, if you look only on the middle of the distribution, it looks the same after clipping away the uninteresting tails, see Figure 7.2; that is, it looks more and more like the normal distribution. This is a universal phenomena known the *central limit theorem* – every sum of nicely behaved random variables behaves like the normal distribution. We unfortunately need a more precise quantification of this behavior, thus the following.

#### 7.1.2. A restricted case of Chernoff inequality via games

##### 7.1.2.1. Chernoff games

**The game.** Consider the game where a player starts with  $Y_0 = 1$  dollars. At every round, the player can bet a certain amount  $x$  (fractions are fine). With probability half she loses her bet, and with probability half she gains an amount equal to her bet. The player is not allowed to go all in – because if she loses then the game is over. So it is natural to ask what her optimal betting strategy is, such that in the end of the game she has as much money as possible.

**Is the game pointless?** So, let  $Y_{i-1}$  be the money the player has in the end of the  $(i-1)$ th round, and she bets an amount  $\psi_i \leq Y_{i-1}$  in the  $i$ th round. As such, in the end of the  $i$ th round, she has

$$Y_i = \begin{cases} Y_{i-1} - \psi_i & \text{LOSE: probability half} \\ Y_{i-1} + \psi_i & \text{WIN: probability half} \end{cases}$$

---

<sup>①</sup>This work is licensed under the Creative Commons Attribution-Noncommercial 3.0 License. To view a copy of this license, visit <http://creativecommons.org/licenses/by-nc/3.0/> or send a letter to Creative Commons, 171 Second Street, Suite 300, San Francisco, California, 94105, USA.

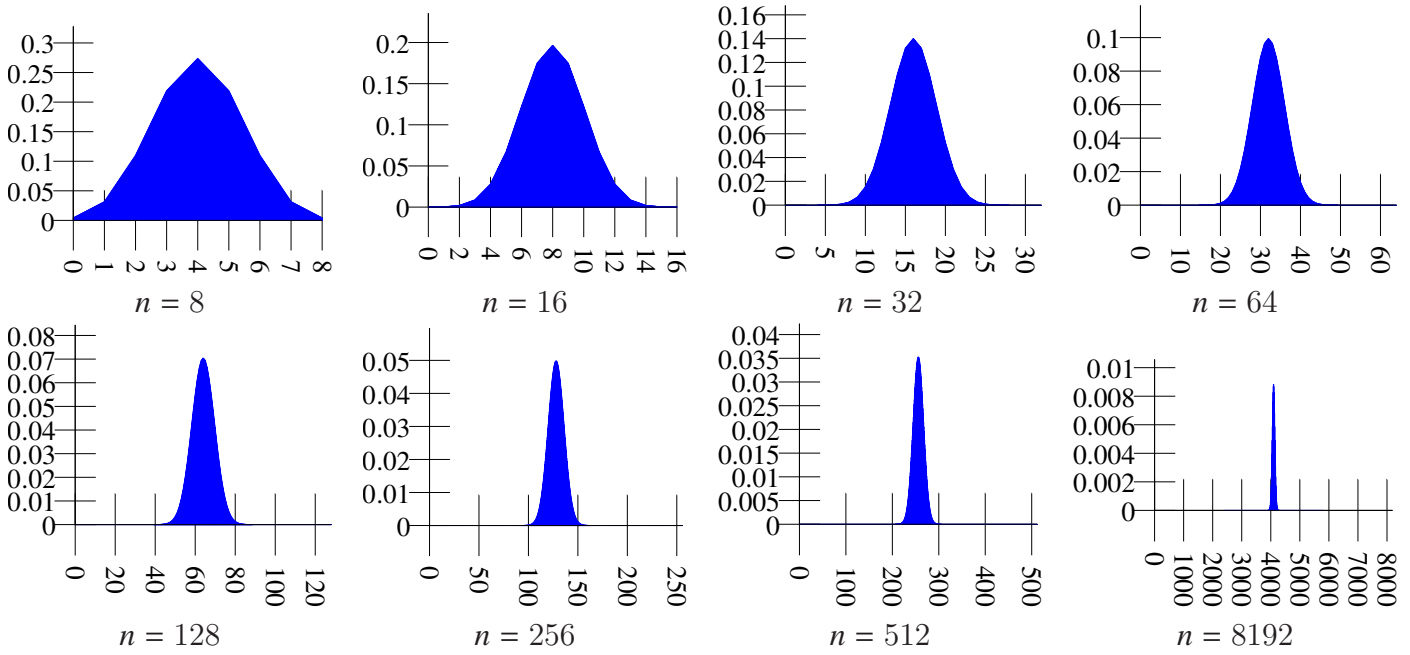


Figure 7.1: The binomial distribution for different values of  $n$ . It pretty quickly concentrates around its expectation.

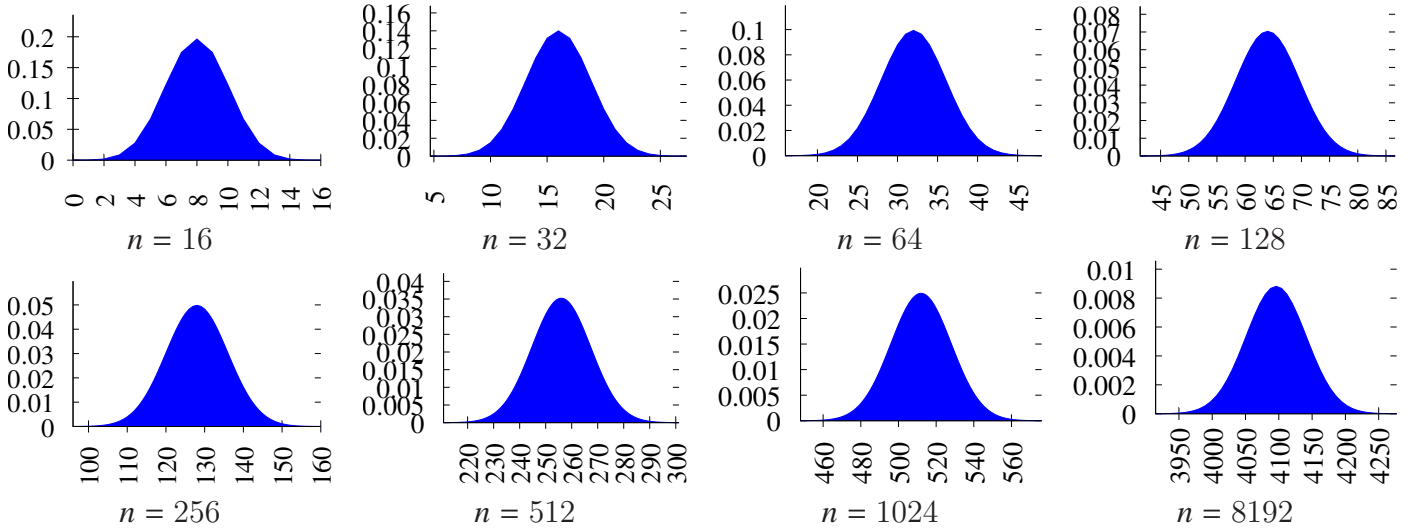


Figure 7.2: The “middle” of the binomial distribution for different values of  $n$ . It very quickly converges to the normal distribution (under appropriate rescaling and translation).

$X_i \in \{-1, +1\}$ $\mathbb{P}[X_i = -1] = \mathbb{P}[X_i = 1] = 1/2$		$X_i \in \{0, 1\}$ $\mathbb{P}[X_i = 0] = \mathbb{P}[X_i = 1] = 1/2$	
$\mathbb{P}[Y \geq \Delta] \leq \exp(-\Delta^2/2n)$	Theorem 7.1.7	$\mathbb{P}[ Y - n/2  \geq \Delta] \leq 2 \exp(-2\Delta^2/n)$	Corollary 7.1.9
$\mathbb{P}[Y \leq -\Delta] \leq \exp(-\Delta^2/2n)$	Theorem 7.1.7		

$X_i \in \{0, 1\}$	$\mathbb{P}[X_i = 1] = p_i$ $\mathbb{P}[X_i = 0] = 1 - p_i$	
$\delta \geq 0$	$P = \mathbb{P}[Y > (1 + \delta)\mu] < \left(e^\delta / (1 + \delta)^{1+\delta}\right)^\mu$	Theorem 7.2.1
$\delta \in (0, 1)$	$P < \exp(-\mu\delta^2/3)$	Lemma 7.2.4
$\delta \in (0, 4)$	$P < \exp(-\mu\delta^2/4)$	Lemma 7.2.5
$\delta \in (0, 6)$	$P < \exp(-\mu\delta^2/5)$	Lemma 7.2.6
$\delta \geq 2e - 1$	$P < 2^{-\mu(1+\delta)}$	Lemma 7.2.7
$\delta \geq e^2$	$P < \exp(-(\mu\delta/2) \ln \delta)$	Lemma 7.2.8
$\delta \geq 0$	$\mathbb{P}[Y < (1 - \delta)\mu] < \left(e^{-\delta} / (1 - \delta)^{1-\delta}\right)^\mu$ $\mathbb{P}[Y < (1 - \delta)\mu] < \exp(-\mu\delta^2/2)$	Theorem 7.2.3
$\Delta \geq 0$	$\mathbb{P}[Y - \mu \geq \Delta] \leq \exp(-2\Delta^2/n)$ $\mathbb{P}[Y - \mu \leq -\Delta] \leq \exp(-2\Delta^2/n)$ .	Corollary 7.3.5
$\tau \geq 1$	$\mathbb{P}[Y < \mu/\tau] < \exp\left(-\left(1 - \frac{1+\ln \tau}{\tau}\right)\mu\right)$	Theorem 7.2.3

$X_i \in [0, 1]$	Arbitrary independent distributions	
$\delta \in [0, 1]$	$\mathbb{P}[Y \geq (1 + \delta)\mu] \leq \exp(-\delta^2\mu/4)$ $\mathbb{P}[Y \leq (1 - \delta)\mu] \leq \exp(-\delta^2\mu/2)$ .	Theorem 7.3.6
$\Delta \geq 0$	$\mathbb{P}[Y - \mu \geq \Delta] \leq \exp(-2\Delta^2/n)$ $\mathbb{P}[Y - \mu \leq -\Delta] \leq \exp(-2\Delta^2/n)$ .	Corollary 7.3.5

$X_i \in [a_i, b_i]$	Arbitrary independent distributions	
$\Delta \geq 0$	$\mathbb{P}[ Y - \mu  \geq \Delta] \leq 2 \exp\left(-\frac{2\Delta^2}{\sum_{i=1}^n (b_i - a_i)^2}\right)$	Theorem 7.4.3

Table 7.1: Summary of Chernoff type inequalities covered. Here we have  $n$  independent random variables  $X_1, \dots, X_n$ ,  $Y = \sum_i X_i$  and  $\mu = \mathbb{E}[Y]$ .

dollars. This game, in expectation, does not change the amount of money the player has. Indeed, we have

$$\mathbb{E}[Y_i | Y_{i-1}] = \frac{1}{2}(Y_{i-1} - \psi_i) + \frac{1}{2}(Y_{i-1} + \psi_i) = Y_{i-1}.$$

And as such, we have that  $\mathbb{E}[Y_i] = \mathbb{E}[\mathbb{E}[Y_i|Y_{i-1}]] = \mathbb{E}[Y_{i-1}] = \dots = \mathbb{E}[Y_0] = 1$ . In particular,  $\mathbb{E}[Y_n] = 1$  – namely, on average, independent of the player strategy she is not going to make any money in this game (and she is allowed to change her bets after every round). Unless, she is lucky<sup>②</sup>...

**What about a lucky player?** The player believes she will get lucky and wants to develop a strategy to take advantage of it. Formally, she believes that she can win, say, at least  $(1 + \delta)/2$  fraction of her bets (instead of the predicted  $1/2$ ) – for example, if the bets are in the stock market, she can improve her chances by doing more research on the companies she is investing in<sup>③</sup>. Unfortunately, the player does not know which rounds she is going to be lucky in – so she still needs to be careful.

**In a search of a good strategy.** Of course, there are many safe strategies the player can use, from not playing at all, to risking only a tiny fraction of her money at each round. In other words, our quest here is to find the best strategy that extracts the maximum benefit for the player out of her inherent luck.

Here, we restrict ourselves to a simple strategy – at every round, the player would bet  $\beta$  fraction of her money, where  $\beta$  is a parameter to be determined. Specifically, in the end of the  $i$ th round, the player would have

$$Y_i = \begin{cases} (1 - \beta)Y_{i-1} & \text{LOSE} \\ (1 + \beta)Y_{i-1} & \text{WIN.} \end{cases}$$

By our assumption, the player is going to win in at least  $M = (1 + \delta)n/2$  rounds. Our purpose here is to figure out what the value of  $\beta$  should be so that player gets as rich as possible<sup>④</sup>. Now, if the player is successful in  $\geq M$  rounds, out of the  $n$  rounds of the game, then the amount of money the player has, in the end of the game, is

$$\begin{aligned} Y_n &\geq (1 - \beta)^{n-M}(1 + \beta)^M = (1 - \beta)^{n/2 - (\delta/2)n}(1 + \beta)^{n/2 + (\delta/2)n} = \left((1 - \beta)(1 + \beta)\right)^{n/2 - (\delta/2)n} (1 + \beta)^{\delta n} \\ &= \left(1 - \beta^2\right)^{n/2 - (\delta/2)n} (1 + \beta)^{\delta n} \geq \exp(-2\beta^2)^{n/2 - (\delta/2)n} \exp(\beta/2)^{\delta n} = \exp((- \beta^2 + \beta^2 \delta + \beta \delta/2)n). \end{aligned}$$

To maximize this quantity, we choose  $\beta = \delta/4$  (there is a better choice, see Lemma 7.1.6, but we use this value for the simplicity of exposition). Thus, we have that  $Y_n \geq \exp\left(\left(-\frac{\delta^2}{16} + \frac{\delta^3}{16} + \frac{\delta^2}{8}\right)n\right) \geq \exp\left(\frac{\delta^2}{16}n\right)$ , proving the following.

**Lemma 7.1.1.** *Consider a Chernoff game with  $n$  rounds, starting with one dollar, where the player wins in  $\geq (1 + \delta)n/2$  of the rounds. If the player bets  $\delta/4$  fraction of her current money, at all rounds, then in the end of the game the player would have at least  $\exp(n\delta^2/16)$  dollars.*

<sup>②</sup>“I would rather have a general who was lucky than one who was good.” – Napoleon Bonaparte.

<sup>③</sup>“I am a great believer in luck, and I find the harder I work, the more I have of it.” – Thomas Jefferson.

<sup>④</sup>This optimal choice is known as Kelly criterion, see Remark 7.1.3.

Remark 7.1.2. Note, that Lemma 7.1.1 holds if the player wins any  $\geq (1 + \delta)n/2$  rounds. In particular, the statement does not require randomness by itself – for our application, however, it is more natural and interesting to think about the player wins as being randomly distributed.

Remark 7.1.3. Interestingly, the idea of choosing the best fraction to bet is an old and natural question arising in investments strategies, and the right fraction to use is known as *Kelly criterion*, going back to Kelly’s work from 1956 [Kel56].

### 7.1.2.2. Chernoff’s inequality

The above implies that if a player is lucky, then she is going to become filthy rich<sup>5</sup>. Intuitively, this should be a pretty rare event – because if the player is rich, then (on average) many other people have to be poor. We are thus ready for the kill.

**Theorem 7.1.4 (Chernoff’s inequality).** *Let  $X_1, \dots, X_n$  be  $n$  independent random variables, where  $X_i = 0$  or  $X_i = 1$  with equal probability. Then, for any  $\delta \in (0, 1/2)$ , we have that*

$$\mathbb{P}\left[\sum_i X_i \geq (1 + \delta)\frac{n}{2}\right] \leq \exp\left(-\frac{\delta^2}{16}n\right).$$

*Proof:* Imagine that we are playing the Chernoff game above, with  $\beta = \delta/4$ , starting with 1 dollar, and let  $Y_i$  be the amount of money in the end of the  $i$ th round. Here  $X_i = 1$  indicates that the player won the  $i$ th round. We have, by Lemma 7.1.1 and Markov’s inequality, that

$$\mathbb{P}\left[\sum_i X_i \geq (1 + \delta)\frac{n}{2}\right] \leq \mathbb{P}\left[Y_n \geq \exp\left(\frac{n\delta^2}{16}\right)\right] \leq \frac{\mathbb{E}[Y_n]}{\exp(n\delta^2/16)} = \frac{1}{\exp(n\delta^2/16)} = \exp\left(-\frac{\delta^2}{16}n\right). \quad \blacksquare$$

**This is crazy – so intuition maybe?** If the player is  $(1 + \delta)/2$ -lucky then she can make a lot of money; specifically, at least  $f(\delta) = \exp(n\delta^2/16)$  dollars by the end of the game. Namely, beating the odds has significant monetary value, and this value grows quickly with  $\delta$ . Since we are in a “zero-sum” game settings, this event should be very rare indeed. Under this interpretation, of course, the player needs to know in advance the value of  $\delta$  – so imagine that she guesses it somehow in advance, or she plays the game in parallel with all the possible values of  $\delta$ , and she settles on the instance that maximizes her profit.

**Can one do better?** No, not really. Chernoff inequality is tight (this is a challenging homework exercise) up to the constant in the exponent. The best bound I know for this version of the inequality has  $1/2$  instead of  $1/16$  in the exponent. Note, however, that no real effort was taken to optimize the constants – this is not the purpose of this write-up.

### 7.1.2.3. Some low level boring calculations

Above, we used the following well known facts.

**Lemma 7.1.5.** (A) *Markov’s inequality.* For any positive random variable  $X$  and  $t > 0$ , we have  $\mathbb{P}[X \geq t] \leq \mathbb{E}[X]/t$ . (B) For any two random variables  $X$  and  $Y$ , we have that  $\mathbb{E}[X] = \mathbb{E}[\mathbb{E}[X|Y]]$ . (C) For  $x \in (0, 1)$ ,  $1 + x \geq e^{x/2}$ . (D) For  $x \in (0, 1/2)$ ,  $1 - x \geq e^{-2x}$ .

<sup>5</sup>Not that there is anything wrong with that – many of my friends are filthy,

**Lemma 7.1.6.** *The quantity  $\exp((-\beta^2 + \beta^2\delta + \beta\delta/2)n)$  is maximal for  $\beta = \frac{\delta}{4(1-\delta)}$ .*

*Proof:* We have to maximize  $f(\beta) = -\beta^2 + \beta^2\delta + \beta\delta/2$  by choosing the correct value of  $\beta$  (as a function of  $\delta$ , naturally).  $f'(\beta) = -2\beta + 2\beta\delta + \delta/2 = 0 \iff 2(\delta - 1)\beta = -\delta/2 \iff \beta = \frac{\delta}{4(1-\delta)}$ . ■

### 7.1.3. A proof for $-1/+1$ case

**Theorem 7.1.7.** *Let  $X_1, \dots, X_n$  be  $n$  independent random variables, such that  $\mathbb{P}[X_i = 1] = \mathbb{P}[X_i = -1] = \frac{1}{2}$ , for  $i = 1, \dots, n$ . Let  $Y = \sum_{i=1}^n X_i$ . Then, for any  $\Delta > 0$ , we have*

$$\mathbb{P}[Y \geq \Delta] \leq \exp(-\Delta^2/2n).$$

*Proof:* Clearly, for an arbitrary  $t$ , to specified shortly, we have

$$\mathbb{P}[Y \geq \Delta] = \mathbb{P}[\exp(tY) \geq \exp(t\Delta)] \leq \frac{\mathbb{E}[\exp(tY)]}{\exp(t\Delta)},$$

the first part follows by the fact that  $\exp(\cdot)$  preserve ordering, and the second part follows by the Markov inequality.

Observe that

$$\begin{aligned} \mathbb{E}[\exp(tX_i)] &= \frac{1}{2}e^t + \frac{1}{2}e^{-t} = \frac{e^t + e^{-t}}{2} \\ &= \frac{1}{2} \left( 1 + \frac{t}{1!} + \frac{t^2}{2!} + \frac{t^3}{3!} + \dots \right) \\ &\quad + \frac{1}{2} \left( 1 - \frac{t}{1!} + \frac{t^2}{2!} - \frac{t^3}{3!} + \dots \right) \\ &= \left( 1 + \frac{t^2}{2!} + \dots + \frac{t^{2k}}{(2k)!} + \dots \right), \end{aligned}$$

by the Taylor expansion of  $\exp(\cdot)$ . Note, that  $(2k)! \geq (k!)2^k$ , and thus

$$\mathbb{E}[\exp(tX_i)] = \sum_{i=0}^{\infty} \frac{t^{2i}}{(2i)!} \leq \sum_{i=0}^{\infty} \frac{t^{2i}}{2^i(i!)} = \sum_{i=0}^{\infty} \frac{1}{i!} \left(\frac{t^2}{2}\right)^i = \exp(t^2/2),$$

again, by the Taylor expansion of  $\exp(\cdot)$ . Next, by the independence of the  $X_i$ s, we have

$$\mathbb{E}[\exp(tY)] = \mathbb{E} \left[ \exp \left( \sum_i tX_i \right) \right] = \mathbb{E} \left[ \prod_i \exp(tX_i) \right] = \prod_{i=1}^n \mathbb{E}[\exp(tX_i)] \leq \prod_{i=1}^n e^{t^2/2} = e^{nt^2/2}.$$

We have  $\mathbb{P}[Y \geq \Delta] \leq \frac{\exp(nt^2/2)}{\exp(t\Delta)} = \exp(nt^2/2 - t\Delta)$ .

Next, by minimizing the above quantity for  $t$ , we set  $t = \Delta/n$ . We conclude,

$$\mathbb{P}[Y \geq \Delta] \leq \exp \left( \frac{n}{2} \left( \frac{\Delta}{n} \right)^2 - \frac{\Delta}{n} \Delta \right) = \exp \left( -\frac{\Delta^2}{2n} \right). \quad \blacksquare$$

By the symmetry of  $Y$ , we get the following:

**Corollary 7.1.8.** Let  $X_1, \dots, X_n$  be  $n$  independent random variables, such that  $\mathbb{P}[X_i = 1] = \mathbb{P}[X_i = -1] = \frac{1}{2}$ , for  $i = 1, \dots, n$ . Let  $Y = \sum_{i=1}^n X_i$ . Then, for any  $\Delta > 0$ , we have  $\mathbb{P}[|Y| \geq \Delta] \leq 2 \exp(-\Delta^2/2n)$ .

**Corollary 7.1.9.** Let  $X_1, \dots, X_n$  be  $n$  independent coin flips, such that  $\mathbb{P}[X_i = 0] = \mathbb{P}[X_i = 1] = \frac{1}{2}$ , for  $i = 1, \dots, n$ . Let  $Y = \sum_{i=1}^n X_i$ . Then, for any  $\Delta > 0$ , we have  $\mathbb{P}[|Y - n/2| \geq \Delta] \leq 2 \exp(-2\Delta^2/n)$ .

**Remark 7.1.10.** Before going any further, it might be instrumental to understand what these inequalities imply. Consider the case where  $X_i$  is either zero or one with probability half. In this case  $\mu = \mathbb{E}[Y] = n/2$ . Set  $\delta = t\sqrt{n}$  ( $\sqrt{\mu}$  is approximately the standard deviation of  $X$  if  $p_i = 1/2$ ). We have by

$$\mathbb{P}\left[\left|Y - \frac{n}{2}\right| \geq \Delta\right] \leq 2 \exp(-2\Delta^2/n) = 2 \exp\left(-2(t\sqrt{n})^2/n\right) = 2 \exp(-2t^2).$$

Thus, Chernoff inequality implies exponential decay (i.e.,  $\leq 2^{-t}$ ) with  $t$  standard deviations, instead of just polynomial (i.e.,  $\leq 1/t^2$ ) by the Chebychev's inequality.

## 7.2. The Chernoff Bound — General Case

Here we present the Chernoff bound in a more general settings.

**Theorem 7.2.1.** Let  $X_1, \dots, X_n$  be  $n$  independent variables, where  $\mathbb{P}[X_i = 1] = p_i$  and  $\mathbb{P}[X_i = 0] = q_i = 1 - p_i$ , for all  $i$ . Let  $X = \sum_{i=1}^n X_i$ .  $\mu = \mathbb{E}[X] = \sum_i p_i$ . For any  $\delta > 0$ , we have

$$\mathbb{P}[X > (1 + \delta)\mu] < \left(e^\delta / (1 + \delta)^{1+\delta}\right)^\mu.$$

*Proof:* We have  $\mathbb{P}[X > (1 + \delta)\mu] = \mathbb{P}\left[e^{tX} > e^{t(1+\delta)\mu}\right]$ . By the Markov inequality, we have:

$$\mathbb{P}[X > (1 + \delta)\mu] < \frac{\mathbb{E}[e^{tX}]}{e^{t(1+\delta)\mu}}$$

On the other hand,

$$\mathbb{E}[e^{tX}] = \mathbb{E}\left[e^{t(X_1 + X_2 + \dots + X_n)}\right] = \mathbb{E}[e^{tX_1}] \dots \mathbb{E}[e^{tX_n}].$$

Namely,

$$\mathbb{P}[X > (1 + \delta)\mu] < \frac{\prod_{i=1}^n \mathbb{E}[e^{tX_i}]}{e^{t(1+\delta)\mu}} = \frac{\prod_{i=1}^n ((1 - p_i)e^0 + p_i e^t)}{e^{t(1+\delta)\mu}} = \frac{\prod_{i=1}^n (1 + p_i(e^t - 1))}{e^{t(1+\delta)\mu}}.$$

Let  $y = p_i(e^t - 1)$ . We know that  $1 + y < e^y$  (since  $y > 0$ ). Thus,

$$\begin{aligned} \mathbb{P}[X > (1 + \delta)\mu] &< \frac{\prod_{i=1}^n \exp(p_i(e^t - 1))}{e^{t(1+\delta)\mu}} = \frac{\exp(\sum_{i=1}^n p_i(e^t - 1))}{e^{t(1+\delta)\mu}} \\ &= \frac{\exp((e^t - 1) \sum_{i=1}^n p_i)}{e^{t(1+\delta)\mu}} = \frac{\exp((e^t - 1)\mu)}{e^{t(1+\delta)\mu}} = \left(\frac{\exp(e^t - 1)}{e^{t(1+\delta)}}\right)^\mu \\ &= \left(\frac{\exp(\delta)}{(1 + \delta)^{1+\delta}}\right)^\mu, \end{aligned}$$

if we set  $t = \log(1 + \delta)$ . ■

### 7.2.1. The lower tail

We need the following low level lemma.

**Lemma 7.2.2.** For  $x \in [0, 1)$ , we have  $(1 - x)^{1-x} \geq \exp(-x + x^2/2)$ .

*Proof:* For  $x \in [0, 1)$ , we have, by the Taylor expansion, that  $\ln(1 - x) = -\sum_{i=1}^{\infty} (x^i/i)$ . As such, we have

$$(1 - x) \ln(1 - x) = -(1 - x) \sum_{i=1}^{\infty} \frac{x^i}{i} = -\sum_{i=1}^{\infty} \frac{x^i}{i} + \sum_{i=1}^{\infty} \frac{x^{i+1}}{i} = -x + \sum_{i=2}^{\infty} \left( \frac{x^i}{i-1} - \frac{x^i}{i} \right) = -x + \sum_{i=2}^{\infty} \frac{x^i}{i(i-1)}.$$

This implies that  $(1 - x) \ln(1 - x) \geq -x + x^2/2$ , which implies the claim by exponentiation.  $\blacksquare$

**Theorem 7.2.3.** Let  $X_1, \dots, X_n$  be  $n$  independent random variables, where  $\mathbb{P}[X_i = 1] = p_i$ ,  $\mathbb{P}[X_i = 0] = q_i = 1 - p_i$ , for all  $i$ . For  $X = \sum_{i=1}^n X_i$ , its expectation is  $\mu = \mathbb{E}[X] = \sum_i p_i$ . We have that

$$\mathbb{P}[X < (1 - \delta)\mu] < \left[ \frac{e^{-\delta}}{(1 - \delta)^{1-\delta}} \right]^{\mu} \quad \text{or alternatively} \quad \mathbb{P}[X < (1 - \delta)\mu] < \exp(-\mu\delta^2/2).$$

For any positive  $\tau > 1$ , we have that  $\mathbb{P}[X < \mu/\tau] \leq \exp\left(-\left(1 - \frac{1+\ln\tau}{\tau}\right)\mu\right)$ .

*Proof:* We follow the same proof template seen already. For  $t = -\ln(1 - \delta) > 0$ , we have  $\mathbb{E}[\exp(-tX_i)] = (1 - p_i)e^0 + p_i e^{-t} = 1 - p_i + p_i(1 - \delta) = 1 - p_i\delta \leq \exp(-p_i\delta)$ . As such, we have

$$\begin{aligned} \mathbb{P}[X < (1 - \delta)\mu] &= \mathbb{P}[-X > -(1 - \delta)\mu] = \mathbb{P}[\exp(-tX) > \exp(-t(1 - \delta)\mu)] \leq \frac{\prod_{i=1}^n \mathbb{E}[\exp(-tX_i)]}{\exp(-t(1 - \delta)\mu)} \\ &\leq \frac{\exp(-\sum_{i=1}^n p_i\delta)}{\exp(-t(1 - \delta)\mu)} = \left[ \frac{e^{-\delta}}{(1 - \delta)^{1-\delta}} \right]^{\mu}. \end{aligned}$$

The alternative simplified form, follows readily from [Lemma 7.2.2](#), since

$$\mathbb{P}[X < (1 - \delta)\mu] \leq \left[ \frac{e^{-\delta}}{(1 - \delta)^{1-\delta}} \right]^{\mu} \leq \left[ \frac{e^{-\delta}}{\exp(-\delta + \delta^2/2)} \right]^{\mu} \leq \exp(-\mu\delta^2/2).$$

For the last inequality, set  $\delta = 1 - 1/\tau$ , and observe that

$$\mathbb{P}[X < (1 - \delta)\mu] \leq \left[ \frac{e^{-\delta}}{(1 - \delta)^{1-\delta}} \right]^{\mu} = \left[ \frac{\exp(-1 + 1/\tau)}{(1/\tau)^{1/\tau}} \right]^{\mu} = \exp\left(-\left(1 - \frac{1 + \ln\tau}{\tau}\right)\mu\right). \quad \blacksquare$$

### 7.2.2. A more convenient form of Chernoff's inequality

**Lemma 7.2.4.** Let  $X_1, \dots, X_n$  be  $n$  independent Bernoulli trials, where  $\mathbb{P}[X_i = 1] = p_i$ , and  $\mathbb{P}[X_i = 0] = 1 - p_i$ , for  $i = 1, \dots, n$ . Let  $X = \sum_{i=1}^n X_i$ , and  $\mu = \mathbb{E}[X] = \sum_i p_i$ . For  $\delta \in (0, 1)$ , we have

$$\mathbb{P}[X > (1 + \delta)\mu] < \exp(-\mu\delta^2/3).$$



*Proof:* By [Theorem 7.2.1](#), it is sufficient to prove, for  $\delta \in [0, 1]$ , that

$$\begin{aligned} \left( \frac{e^\delta}{(1+\delta)^{1+\delta}} \right)^\mu \leq \exp\left(-\frac{\mu\delta^2}{c}\right) &\iff \mu(\delta - (1+\delta)\ln(1+\delta)) \leq -\mu\delta^2/c \\ \iff f(\delta) = \delta^2/c + \delta - (1+\delta)\ln(1+\delta) &\leq 0. \end{aligned}$$

We have

$$f'(\delta) = 2\delta/c - \ln(1+\delta). \quad \text{and} \quad f''(\delta) = 2/c - \frac{1}{1+\delta}.$$

For  $c = 3$ , we have  $f''(\delta) \leq 0$  for  $\delta \in [0, 1/2]$ , and  $f''(\delta) \geq 0$  for  $\delta \in [1/2, 1]$ . Namely,  $f'(\delta)$  achieves its maximum either at 0 or 1. As  $f'(0) = 0$  and  $f'(1) = 2/3 - \ln 2 \approx -0.02 < 0$ , we conclude that  $f'(\delta) \leq 0$ . Namely,  $f$  is a monotonically decreasing function in  $[0, 1]$ , which implies that  $f(\delta) \leq 0$ , for all  $\delta$  in this range, thus implying the claim.  $\blacksquare$

**Lemma 7.2.5.** *Let  $X_1, \dots, X_n$  be  $n$  independent Bernoulli trials, where  $\mathbb{P}[X_i = 1] = p_i$ , and  $\mathbb{P}[X_i = 0] = 1 - p_i$ , for  $i = 1, \dots, n$ . Let  $X = \sum_{i=1}^n X_i$ , and  $\mu = \mathbb{E}[X] = \sum_i p_i$ . For  $\delta \in (0, 4)$ , we have*

$$\mathbb{P}[X > (1+\delta)\mu] < \exp(-\mu\delta^2/4),$$

*Proof:* [Lemma 7.2.4](#) implies a stronger bound, so we need to prove the claim only for  $\delta \in (1, 4]$ . Continuing as in the proof of [Lemma 7.2.4](#), for case  $c = 4$ , we have to prove that

$$f(\delta) = \delta^2/4 + \delta - (1+\delta)\ln(1+\delta) \leq 0,$$

where  $f''(\delta) = 1/2 - \frac{1}{1+\delta}$ .

For  $\delta > 1$ , we have  $f''(\delta) > 0$ . Namely  $f(\cdot)$  is convex for  $\delta \geq 1$ , and it achieves its maximum on the interval  $[1, 4]$  on the endpoints. In particular,  $f(1) \approx -0.13$ , and  $f(4) \approx -0.047$ , which implies the claim.  $\blacksquare$

**Lemma 7.2.6.** *Let  $X_1, \dots, X_n$  be  $n$  independent random variables, where  $\mathbb{P}[X_i = 1] = p_i$ , and  $\mathbb{P}[X_i = 0] = 1 - p_i$ , for  $i = 1, \dots, n$ . Let  $X = \sum_{i=1}^n X_i$ , and  $\mu = \mathbb{E}[X] = \sum_i p_i$ . For  $\delta \in (0, 6)$ , we have*

$$\mathbb{P}[X > (1+\delta)\mu] < \exp(-\mu\delta^2/5),$$

*Proof:* [Lemma 7.2.5](#) implies a stronger bound, so we need to prove the claim only for  $\delta \in (4, 5]$ . Continuing as in the proof of [Lemma 7.2.4](#), for case  $c = 5$ , we have to prove that

$$f(\delta) = \delta^2/5 + \delta - (1+\delta)\ln(1+\delta) \leq 0,$$

where  $f''(\delta) = 2/5 - \frac{1}{1+\delta}$ . For  $\delta \geq 4$ , we have  $f''(\delta) > 0$ . Namely  $f(\cdot)$  is convex for  $\delta \geq 4$ , and it achieves its maximum on the interval  $[4, 6]$  on the endpoints. In particular,  $f(4) \approx -0.84$ , and  $f(6) \approx -0.42$ , which implies the claim.  $\blacksquare$

**Lemma 7.2.7.** *Let  $X_1, \dots, X_n$  be  $n$  independent Bernoulli trials, where  $\mathbb{P}[X_i = 1] = p_i$ , and  $\mathbb{P}[X_i = 0] = 1 - p_i$ , for  $i = 1, \dots, n$ . Let  $X = \sum_{i=1}^n X_i$ , and  $\mu = \mathbb{E}[X] = \sum_i p_i$ . For  $\delta > 2e - 1$ , we have  $\mathbb{P}[X > (1+\delta)\mu] < 2^{-\mu(1+\delta)}$ .*

*Proof:* By Theorem 7.2.1, we have

$$\left(\frac{e}{1+\delta}\right)^{(1+\delta)\mu} \leq \left(\frac{e}{1+2e-1}\right)^{(1+\delta)\mu} \leq 2^{-(1+\delta)\mu},$$

since  $\delta > 2e - 1$ . ■

**Lemma 7.2.8.** *Let  $X_1, \dots, X_n$  be  $n$  independent Bernoulli trials, where  $\mathbb{P}[X_i = 1] = p_i$ , and  $\mathbb{P}[X_i = 0] = 1 - p_i$ , for  $i = 1, \dots, n$ . Let  $X = \sum_{i=1}^n X_i$ , and  $\mu = \mathbb{E}[X] = \sum_i p_i$ . For  $\delta > e^2$ , we have  $\mathbb{P}[X > (1 + \delta)\mu] < \exp\left(-\frac{\mu\delta \ln \delta}{2}\right)$ .*

*Proof:* Observe that

$$\mathbb{P}[X > (1 + \delta)\mu] < \left(\frac{e^\delta}{(1 + \delta)^{1+\delta}}\right)^\mu = \exp\left(\mu\delta - \mu(1 + \delta) \ln(1 + \delta)\right). \quad (7.1)$$

As such, we have

$$\mathbb{P}[X > (1 + \delta)\mu] < \exp\left(-\mu(1 + \delta)(\ln(1 + \delta) - 1)\right) \leq \exp\left(-\mu\delta \ln \frac{1 + \delta}{e}\right) \leq \exp\left(-\frac{\mu\delta \ln \delta}{2}\right),$$

since for  $x \geq e^2$  we have that  $\frac{1+x}{e} \geq \sqrt{x} \iff \ln \frac{1+x}{e} \geq \frac{\ln x}{2}$ . ■

### 7.3. A special case of Hoeffding's inequality

In this section, we prove yet another version of Chernoff inequality, where each variable is randomly picked according to its own distribution in the range  $[0, 1]$ . We prove a more general version of this inequality in Section 7.4, but the version presented here does not follow from this generalization.

**Theorem 7.3.1.** *Let  $X_1, \dots, X_n \in [0, 1]$  be  $n$  independent random variables, let  $X = \sum_{i=1}^n X_i$ , and let  $\mu = \mathbb{E}[X]$ . We have that  $\mathbb{P}[X - \mu \geq \eta] \leq \left(\frac{\mu}{\mu + \eta}\right)^{\mu + \eta} \left(\frac{n - \mu}{n - \mu - \eta}\right)^{n - \mu - \eta}$ .*

*Proof:* Let  $s \geq 1$  be some arbitrary parameter. By the standard arguments, we have

$$\gamma = \mathbb{P}[X \geq \mu + \eta] = \mathbb{P}[s^X \geq s^{\mu + \eta}] \leq \frac{\mathbb{E}[s^X]}{s^{\mu + \eta}} = s^{-\mu - \eta} \prod_{i=1}^n \mathbb{E}[s^{X_i}].$$

By calculations, see Lemma 7.3.7 below, one can show that  $\mathbb{E}[s^{X_i}] \leq 1 + (s - 1)\mathbb{E}[X_i]$ . As such, by the AM-GM inequality<sup>®</sup>, we have that

$$\prod_{i=1}^n \mathbb{E}[s^{X_i}] \leq \prod_{i=1}^n \left(1 + (s - 1)\mathbb{E}[X_i]\right) \leq \left(\frac{1}{n} \sum_{i=1}^n \left(1 + (s - 1)\mathbb{E}[X_i]\right)\right)^n = \left(1 + (s - 1)\frac{\mu}{n}\right)^n.$$

---

<sup>®</sup>The inequality between arithmetic and geometric means:  $(\sum_{i=1}^n x_i)/n \geq \sqrt[n]{x_1 \cdots x_n}$ .

Setting  $s = \frac{(\mu + \eta)(n - \mu)}{\mu(n - \mu - \eta)} = \frac{\mu n - \mu^2 + \eta n - \eta \mu}{\mu n - \mu^2 - \eta \mu}$  we have that

$$1 + (s - 1)\frac{\mu}{n} = 1 + \frac{\eta n}{\mu n - \mu^2 - \eta \mu} \cdot \frac{\mu}{n} = 1 + \frac{\eta}{n - \mu - \eta} = \frac{n - \mu}{n - \mu - \eta}.$$

As such, we have that

$$\gamma \leq s^{-\mu - \eta} \prod_{i=1}^n \mathbb{E}[s^{X_i}] = \left( \frac{\mu(n - \mu - \eta)}{(\mu + \eta)(n - \mu)} \right)^{\mu + \eta} \left( \frac{n - \mu}{n - \mu - \eta} \right)^n = \left( \frac{\mu}{\mu + \eta} \right)^{\mu + \eta} \left( \frac{n - \mu}{n - \mu - \eta} \right)^{n - \mu - \eta}. \quad \blacksquare$$

**Remark 7.3.2.** Setting  $s = (\mu + \eta)/\mu$  in the proof of [Theorem 7.3.1](#), we have

$$\mathbb{P}\left[X - \mu \geq \eta\right] \leq \left(\frac{\mu}{\mu + \eta}\right)^{\mu + \eta} \left(1 + \left(\frac{\mu + \eta}{\mu} - 1\right)\frac{\mu}{n}\right)^n = \left(\frac{\mu}{\mu + \eta}\right)^{\mu + \eta} \left(1 + \frac{\eta}{n}\right)^n.$$

**Corollary 7.3.3.** Let  $X_1, \dots, X_n \in [0, 1]$  be  $n$  independent random variables, let  $\bar{X} = \sum_{i=1}^n X_i/n$ ,  $p = \mathbb{E}[\bar{X}] = \mu/n$  and  $q = 1 - p$ . Then, we have that  $\mathbb{P}\left[\bar{X} - p \geq t\right] \leq \exp(nf(t))$ , for

$$f(t) = (p + t) \ln \frac{p}{p + t} + (q - t) \ln \frac{q}{q - t}. \quad (7.2)$$

**Theorem 7.3.4.** Let  $X_1, \dots, X_n \in [0, 1]$  be  $n$  independent random variables, let  $\bar{X} = (\sum_{i=1}^n X_i)/n$ , and let  $p = \mathbb{E}[X]$ . We have that  $\mathbb{P}\left[\bar{X} - p \geq t\right] \leq \exp(-2nt^2)$  and  $\mathbb{P}\left[\bar{X} - p \leq -t\right] \leq \exp(-2nt^2)$ .

*Proof:* Let  $p = \mu/n$ ,  $q = 1 - p$ , and let  $f(t)$  be the function from [Eq. \(7.2\)](#), for  $t \in (-p, q)$ . Now, we have that

$$\begin{aligned} f'(t) &= \ln \frac{p}{p + t} + (p + t) \frac{p + t}{p} \left( -\frac{p}{(p + t)^2} \right) - \ln \frac{q}{q - t} - (q - t) \frac{q - t}{q} \frac{q}{(q - t)^2} = \ln \frac{p}{p + t} - \ln \frac{q}{q - t} \\ &= \ln \frac{p(q - t)}{q(p + t)}. \end{aligned}$$

As for the second derivative, we have

$$f''(t) = \frac{q(p + t)}{p(q - t)} \cdot \frac{p}{q} \cdot \frac{(p + t)(-1) - (q - t)}{(p + t)^2} = \frac{-p - t - q + t}{(q - t)(p + t)} = -\frac{1}{(q - t)(p + t)} \leq -4.$$

Indeed,  $t \in (-p, q)$  and the denominator is minimized for  $t = (q - p)/2$ , and as such  $(q - t)(p + t) \leq (2q - (q - p))(2p + (q - p))/4 = (p + q)^2/4 = 1/4$ .

Now,  $f(0) = 0$  and  $f'(0) = 0$ , and by Taylor's expansion, we have that  $f(t) = f(0) + f'(0)t + \frac{f''(x)}{2}t^2 \leq -2t^2$ , where  $x$  is between 0 and  $t$ .

The first bound now readily follows from plugging this bound into [Corollary 7.3.3](#). The second bound follows by considering the random variants  $Y_i = 1 - X_i$ , for all  $i$ , and plugging this into the first bound. Indeed, for  $\bar{Y} = 1 - \bar{X}$ , we have that  $q = \mathbb{E}[\bar{Y}]$ , and then  $\bar{X} - p \leq -t \iff t \leq p - \bar{X} \iff t \leq 1 - q - (1 - \bar{Y}) = \bar{Y} - q$ . Thus,  $\mathbb{P}\left[\bar{X} - p \leq -t\right] = \mathbb{P}\left[\bar{Y} - q \geq t\right] \leq \exp(-2nt^2)$ .  $\blacksquare$

**Corollary 7.3.5.** Let  $X_1, \dots, X_n \in [0, 1]$  be  $n$  independent random variables, let  $Y = \sum_{i=1}^n X_i$ , and let  $\mu = \mathbb{E}[X]$ . For any  $\Delta > 0$ , we have  $\mathbb{P}[Y - \mu \geq \Delta] \leq \exp(-2\Delta^2/n)$  and  $\mathbb{P}[Y - \mu \leq -\Delta] \leq \exp(-2\Delta^2/n)$ .

*Proof:* For  $\bar{X} = Y/n$ ,  $p = \mu/n$ , and  $t = \Delta/n$ , by [Theorem 7.3.4](#), we have

$$\mathbb{P}[Y - \mu \geq \Delta] = \mathbb{P}[\bar{X} - p \geq t] \leq \exp(-2nt^2) = \exp(-2\Delta^2/n). \quad \blacksquare$$

**Theorem 7.3.6.** Let  $X_1, \dots, X_n \in [0, 1]$  be  $n$  independent random variables, let  $X = (\sum_{i=1}^n X_i)$ , and let  $\mu = \mathbb{E}[X]$ . We have that  $\mathbb{P}[X - \mu \geq \varepsilon\mu] \leq \exp(-\varepsilon^2\mu/4)$  and  $\mathbb{P}[X - \mu \leq -\varepsilon\mu] \leq \exp(-\varepsilon^2\mu/2)$ .

*Proof:* Let  $p = \mu/n$ , and let  $g(x) = f(px)$ , for  $x \in [0, 1]$  and  $xp < q$ . As before, computing the derivative of  $g$ , we have

$$g'(x) = pf'(xp) = p \ln \frac{p(q-xp)}{q(p+xp)} = p \ln \frac{q-xp}{q(1+x)} \leq p \ln \frac{1}{1+x} \leq -\frac{px}{2},$$

since  $(q-xp)/q$  is maximized for  $x = 0$ , and  $\ln \frac{1}{1+x} \leq -x/2$ , for  $x \in [0, 1]$ , as can be easily verified<sup>Ⓣ</sup>. Now,  $g(0) = f(0) = 0$ , and by integration, we have that  $g(x) = \int_{y=0}^x g'(y)dy \leq \int_{y=0}^x (-py/2)dy = -px^2/4$ . Now, plugging into [Corollary 7.3.3](#), we get that the desired probability  $\mathbb{P}[X - \mu \geq \varepsilon\mu]$  is

$$\mathbb{P}[\bar{X} - p \geq \varepsilon p] \leq \exp(nf(\varepsilon p)) = \exp(ng(\varepsilon)) \leq \exp(-pn\varepsilon^2/4) = \exp(-\mu\varepsilon^2/4).$$

As for the other inequality, set  $h(x) = g(-x) = f(-xp)$ . Then

$$\begin{aligned} h'(x) &= -pf'(-xp) = -p \ln \frac{p(q+xp)}{q(p-xp)} = p \ln \frac{q(1-x)}{q+xp} = p \ln \frac{q-xq}{q+xp} = p \ln \left(1 - x \frac{p+q}{q+xp}\right) \\ &= p \ln \left(1 - x \frac{1}{q+xp}\right) \leq p \ln(1-x) \leq -px, \end{aligned}$$

since  $1-x \leq e^{-x}$ . By integration, as before, we conclude that  $h(x) \leq -px^2/2$ . Now, plugging into [Corollary 7.3.3](#), we get  $\mathbb{P}[X - \mu \leq -\varepsilon\mu] = \mathbb{P}[\bar{X} - p \leq -\varepsilon p] \leq \exp(nf(-\varepsilon p)) \leq \exp(nh(\varepsilon)) \leq \exp(-np\varepsilon^2/2) \leq \exp(-\mu\varepsilon^2/2)$ .  $\blacksquare$

### 7.3.1. Some technical lemmas

**Lemma 7.3.7.** Let  $X \in [0, 1]$  be a random variable, and let  $s \geq 1$ . Then  $\mathbb{E}[s^X] \leq 1 + (s-1)\mathbb{E}[X]$ .

*Proof:* For the sake of simplicity of exposition, assume that  $X$  is a discrete random variable, and that there is a value  $\alpha \in (0, 1/2)$ , such that  $\beta = \mathbb{P}[X = \alpha] > 0$ . Consider the modified random variable  $X'$ , such that  $\mathbb{P}[X' = 0] = \mathbb{P}[X = 0] + \beta/2$ , and  $\mathbb{P}[X' = 2\alpha] = \mathbb{P}[X = \alpha] + \beta/2$ . Clearly,  $\mathbb{E}[X] = \mathbb{E}[X']$ . Next, observe that  $\mathbb{E}[s^{X'}] - \mathbb{E}[s^X] = (\beta/2)(s^{2\alpha} + s^0) - \beta s^\alpha \geq 0$ , by the convexity of  $s^x$ . We conclude that  $\mathbb{E}[s^X]$  achieves its maximum if it takes only the values 0 and 1. But then, we have that  $\mathbb{E}[s^X] = \mathbb{P}[X = 0]s^0 + \mathbb{P}[X = 1]s^1 = (1 - \mathbb{E}[X]) + \mathbb{E}[X]s = 1 + (s-1)\mathbb{E}[X]$ , as claimed.  $\blacksquare$

<sup>Ⓣ</sup>Indeed, this is equivalent to  $\frac{1}{1+x} \leq e^{-x/2} \iff e^{x/2} \leq 1+x$ , which readily holds for  $x \in [0, 1]$ .

## 7.4. Hoeffding's inequality

In this section, we prove a generalization of Chernoff's inequality. The proof is considerably more tedious, and it is included here for the sake of completeness.

**Lemma 7.4.1.** *Let  $X$  be a random variable. If  $\mathbb{E}[X] = 0$  and  $a \leq X \leq b$ , then for any  $s > 0$ , we have  $\mathbb{E}[e^{sX}] \leq \exp(s^2(b-a)^2/8)$ .*

*Proof:* Let  $a \leq x \leq b$  and observe that  $x$  can be written as a convex combination of  $a$  and  $b$ . In particular, we have

$$x = \lambda a + (1 - \lambda)b \quad \text{for} \quad \lambda = \frac{b - x}{b - a} \in [0, 1].$$

Since  $s > 0$ , the function  $\exp(sx)$  is convex, and as such

$$e^{sx} \leq \frac{b - x}{b - a} e^{sa} + \frac{x - a}{b - a} e^{sb},$$

since we have that  $f(\lambda x + (1 - \lambda)y) \leq \lambda f(x) + (1 - \lambda)f(y)$  if  $f(\cdot)$  is a convex function. Thus, for a random variable  $X$ , by linearity of expectation, we have

$$\begin{aligned} \mathbb{E}[e^{sX}] &\leq \mathbb{E}\left[\frac{b - X}{b - a} e^{sa} + \frac{X - a}{b - a} e^{sb}\right] = \frac{b - \mathbb{E}[X]}{b - a} e^{sa} + \frac{\mathbb{E}[X] - a}{b - a} e^{sb} \\ &= \frac{b}{b - a} e^{sa} - \frac{a}{b - a} e^{sb}, \end{aligned}$$

since  $\mathbb{E}[X] = 0$ .

Next, set  $p = -\frac{a}{b - a}$  and observe that  $1 - p = 1 + \frac{a}{b - a} = \frac{b}{b - a}$  and

$$-ps(b - a) = -\left(-\frac{a}{b - a}\right)s(b - a) = sa.$$

As such, we have

$$\begin{aligned} \mathbb{E}[e^{sX}] &\leq (1 - p)e^{sa} + pe^{sb} = (1 - p + pe^{s(b-a)})e^{sa} \\ &= (1 - p + pe^{s(b-a)})e^{-ps(b-a)} \\ &= \exp\left(-ps(b - a) + \ln(1 - p + pe^{s(b-a)})\right) = \exp(-pu + \ln(1 - p + pe^u)), \end{aligned}$$

for  $u = s(b - a)$ . Setting

$$\phi(u) = -pu + \ln(1 - p + pe^u),$$

we thus have  $\mathbb{E}[e^{sX}] \leq \exp(\phi(u))$ . To prove the claim, we will show that  $\phi(u) \leq u^2/8 = s^2(b - a)^2/8$ .

To see that, expand  $\phi(u)$  about zero using Taylor's expansion. We have

$$\phi(u) = \phi(0) + u\phi'(0) + \frac{1}{2}u^2\phi''(\theta) \tag{7.3}$$

where  $\theta \in [0, u]$ , and notice that  $\phi(0) = 0$ . Furthermore, we have

$$\phi'(u) = -p + \frac{pe^u}{1 - p + pe^u},$$

and as such  $\phi'(0) = -p + \frac{p}{1-p+p} = 0$ . Now,

$$\phi''(u) = \frac{(1-p+pe^u)pe^u - (pe^u)^2}{(1-p+pe^u)^2} = \frac{(1-p)pe^u}{(1-p+pe^u)^2}.$$

For any  $x, y \geq 0$ , we have  $(x+y)^2 \geq 4xy$  as this is equivalent to  $(x-y)^2 \geq 0$ . Setting  $x = 1-p$  and  $y = pe^u$ , we have that

$$\phi''(u) = \frac{(1-p)pe^u}{(1-p+pe^u)^2} \leq \frac{(1-p)pe^u}{4(1-p)pe^u} = \frac{1}{4}.$$

Plugging this into Eq. (7.3), we get that

$$\phi(u) \leq \frac{1}{8}u^2 = \frac{1}{8}(s(b-a))^2 \quad \text{and} \quad \mathbb{E}[e^{sX}] \leq \exp(\phi(u)) \leq \exp\left(\frac{1}{8}(s(b-a))^2\right),$$

as claimed. ■

**Lemma 7.4.2.** *Let  $X$  be a random variable. If  $\mathbb{E}[X] = 0$  and  $a \leq X \leq b$ , then for any  $s > 0$ , we have*

$$\mathbb{P}[X > t] \leq \frac{\exp\left(\frac{s^2(b-a)^2}{8}\right)}{e^{st}}.$$

*Proof:* Using the same technique we used in proving Chernoff's inequality, we have that

$$\mathbb{P}[X > t] = \mathbb{P}[e^{sX} > e^{st}] \leq \frac{\mathbb{E}[e^{sX}]}{e^{st}} \leq \frac{\exp\left(\frac{s^2(b-a)^2}{8}\right)}{e^{st}}. \quad \blacksquare$$

**Theorem 7.4.3 (Hoeffding's inequality).** *Let  $X_1, \dots, X_n$  be independent random variables, where  $X_i \in [a_i, b_i]$ , for  $i = 1, \dots, n$ . Then, for the random variable  $S = X_1 + \dots + X_n$  and any  $\eta > 0$ , we have*

$$\mathbb{P}\left[|S - \mathbb{E}[S]| \geq \eta\right] \leq 2 \exp\left(-\frac{2\eta^2}{\sum_{i=1}^n (b_i - a_i)^2}\right).$$

*Proof:* Let  $Z_i = X_i - \mathbb{E}[X_i]$ , for  $i = 1, \dots, n$ . Set  $Z = \sum_{i=1}^n Z_i$ , and observe that

$$\mathbb{P}[Z \geq \eta] = \mathbb{P}[e^{sZ} \geq e^{s\eta}] \leq \frac{\mathbb{E}[\exp(sZ)]}{\exp(s\eta)},$$

by Markov's inequality. Arguing as in the proof of Chernoff's inequality, we have

$$\mathbb{E}[\exp(sZ)] = \mathbb{E}\left[\prod_{i=1}^n \exp(sZ_i)\right] = \prod_{i=1}^n \mathbb{E}[\exp(sZ_i)] \leq \prod_{i=1}^n \exp\left(\frac{s^2(b_i - a_i)^2}{8}\right),$$

since the  $Z_i$ s are independent and by Lemma 7.4.1. This implies that

$$\mathbb{P}[Z \geq \eta] \leq \exp(-s\eta) \prod_{i=1}^n e^{s^2(b_i - a_i)^2/8} = \exp\left(\frac{s^2}{8} \sum_{i=1}^n (b_i - a_i)^2 - s\eta\right).$$

The upper bound is minimized for  $s = 4\eta / (\sum_i (b_i - a_i)^2)$ , implying

$$\mathbb{P}[Z \geq \eta] \leq \exp\left(-\frac{2\eta^2}{\sum (b_i - a_i)^2}\right).$$

The claim now follows by the symmetry of the upper bound (i.e., apply the same proof to  $-Z$ ). ■

## 7.5. Bibliographical notes

Some of the exposition here follows more or less the exposition in [MR95]. Exercise 7.6.1 (without the hint) is from [Mat99]. McDiarmid [McD89] provides a survey of Chernoff type inequalities, and Theorem 7.3.6 and Section 7.3 is taken from there (our proof has somewhat weaker constants).

A more general treatment of such inequalities and tools is provided by Dubhashi and Panconesi [DP09].

## 7.6. Exercises

Exercise 7.6.1 (Chernoff inequality is tight.). Let  $S = \sum_{i=1}^n S_i$  be a sum of  $n$  independent random variables each attaining values  $+1$  and  $-1$  with equal probability. Let  $P(n, \Delta) = \mathbb{P}[S > \Delta]$ . Prove that for  $\Delta \leq n/C$ ,

$$P(n, \Delta) \geq \frac{1}{C} \exp\left(-\frac{\Delta^2}{Cn}\right),$$

where  $C$  is a suitable constant. That is, the well-known Chernoff bound  $P(n, \Delta) \leq \exp(-\Delta^2/2n)$  is close to the truth.

Exercise 7.6.2 (Chernoff inequality is tight by direct calculations.). For this question use only basic argumentation – do not use Stirling’s formula, Chernoff inequality or any similar “heavy” machinery.

(A) Prove that 
$$\sum_{i=0}^{n-k} \binom{2n}{i} \leq \frac{n}{4k^2} 2^{2n}.$$

Hint: Consider flipping a coin  $2n$  times. Write down explicitly the probability of this coin to have at most  $n - k$  heads, and use Chebyshev inequality.

(B) Using (A), prove that  $\binom{2n}{n} \geq 2^{2n}/4\sqrt{n}$  (which is a pretty good estimate).

(C) Prove that 
$$\binom{2n}{n+i+1} = \left(1 - \frac{2i+1}{n+i+1}\right) \binom{2n}{n+i}.$$

(D) Prove that 
$$\binom{2n}{n+i} \leq \exp\left(\frac{-i(i-1)}{2n}\right) \binom{2n}{n}.$$

(E) Prove that 
$$\binom{2n}{n+i} \geq \exp\left(-\frac{8i^2}{n}\right) \binom{2n}{n}.$$

(F) Using the above, prove that  $\binom{2n}{n} \leq c \frac{2^{2n}}{\sqrt{n}}$  for some constant  $c$  (I got  $c = 0.824\dots$  but any reasonable constant will do).

(G) Using the above, prove that

$$\sum_{i=t\sqrt{n}+1}^{(t+1)\sqrt{n}} \binom{2n}{n-i} \leq c 2^{2n} \exp(-t^2/2).$$

In particular, conclude that when flipping fair coin  $2n$  times, the probability to get less than  $n - t\sqrt{n}$  heads (for  $t$  an integer) is smaller than  $c' \exp(-t^2/2)$ , for some constant  $c'$ .

(H) Let  $X$  be the number of heads in  $2n$  coin flips. Prove that for any integer  $t > 0$  and any  $\delta > 0$  sufficiently small, it holds that  $\mathbb{P}[X < (1 - \delta)n] \geq \exp(-c''\delta^2n)$ , where  $c''$  is some constant. Namely, the Chernoff inequality is tight in the worst case.

Exercise 7.6.3 (Tail inequality for geometric variables). Let  $X_1, \dots, X_m$  be  $m$  independent random variables with geometric distribution with probability  $p$  (i.e.,  $\mathbb{P}[X_i = j] = (1 - p)^{j-1}p$ ). Let  $Y = \sum_i X_i$ , and let  $\mu = \mathbb{E}[Y] = m/p$ . Prove that  $\mathbb{P}[Y \geq (1 + \delta)\mu] \leq \exp(-m\delta^2/8)$ .

## Bibliography

- [DP09] Devdatt P. Dubhashi and Alessandro Panconesi. *Concentration of Measure for the Analysis of Randomized Algorithms*. Cambridge University Press, 2009.
- [Kel56] J. L. Kelly. *A new interpretation of information rate*. *Bell Sys. Tech. J.*, 35(4):917–926, jul 1956.
- [Mat99] J. Matoušek. *Geometric Discrepancy*. Springer, 1999.
- [McD89] C. McDiarmid. *Surveys in Combinatorics*, chapter On the method of bounded differences. Cambridge University Press, 1989.
- [MR95] R. Motwani and P. Raghavan. *Randomized Algorithms*. Cambridge University Press, Cambridge, UK, 1995.