

The One-Round Voronoi Game*

OTFRIED CHEONG

Institute of Information and Computing Sciences
Utrecht University
P.O.Box 80.089, 3508 TB Utrecht
the Netherlands

SARIEL HAR-PELED

Department of Computer Science
University of Illinois
1304 W. Springfield Ave.
Urbana, IL 61801; USA

NATHAN LINIAL

Institute of Computer Science
Hebrew University
Jerusalem 91904, Israel

JIRÍ MATOUŠEK**

Department of Applied Mathematics and
Institute of Theoretical Computer Science (ITI)
Charles University
Malostranské nám. 25, 118 00 Praha 1
Czech Republic

Abstract

In the one-round Voronoi game, the \mathcal{FRST} player places n sites inside a unit-square Q . Next, the \mathcal{SCND} player places n points inside Q . The payoff for a player is the total area of the Voronoi region of Q under their control. In this paper, we show that the \mathcal{SCND} player can always place the points in such a way that it controls $1/2 + \alpha$ fraction of the total area of Q , where $\alpha > 0$ is a constant independent of n .

1 Introduction

Competitive facility location studies the placement of sites by competing market players. Overviews of different models are the surveys by Friesz et al. [TFM89], Eiselt and Laporte [EL89] and Eiselt et al. [ELT93].

The *Voronoi game* is a simple geometric model for competitive facility location, where a site s “owns” that part of the playing arena that is closer to s than to any other site. We consider a two-player version with a square arena Q . The two players, \mathcal{FRST} and \mathcal{SCND} , place points into Q . The goal of both players is to capture as much of the area of Q as possible, where the region captured by \mathcal{FRST} is $R(\mathcal{F}, \mathcal{S}) = \{x \in Q : \text{dist}(x, \mathcal{F}) < \text{dist}(x, \mathcal{S})\}$ and the region captured by \mathcal{SCND} is $R(\mathcal{S}, \mathcal{F})$. Here \mathcal{F} is the set of points of \mathcal{FRST} , \mathcal{S} is the set of points of \mathcal{SCND} , $\text{dist}(\cdot, \cdot)$ is the Euclidean distance and $\text{vol}(\cdot)$ is the Lebesgue measure. In other words, if we construct the Voronoi diagram [Aur91] of $\mathcal{F} \cup \mathcal{S}$, then each player captures the Voronoi regions (restricted to Q) of his

*Part of this research was done during a workshop supported by ITI (project LN00A056 of the Ministry of Education of the Czech Republic).

**Supported by Charles University grants No. 158/99 and 159/99.

point set and is rewarded proportionate to the measure of his captured set. The *payoff* for \mathcal{FRST} is $\text{vol}(R(\mathcal{F}, \mathcal{S}))/\text{vol}(Q)$ and the payoff for \mathcal{SCND} is $\text{vol}(R(\mathcal{S}, \mathcal{F}))/\text{vol}(Q)$. (Of course, we can re-scale the board Q so that $\text{vol}(Q) = 1$, but in the subsequent considerations a different scaling seems more intuitive.)

Ahn et al. [ACC⁺01] had studied a one-dimensional Voronoi game, where the arena Q is a line segment, and the game takes n rounds. In each round, \mathcal{FRST} and \mathcal{SCND} place one point each. Ahn et al. showed that \mathcal{SCND} then has a winning strategy that guarantees a payoff of $1/2 + \epsilon$, with $\epsilon > 0$, but that \mathcal{FRST} can force ϵ to be as small as he wishes. On the other hand, if only a single round is played, where \mathcal{FRST} first places n points, followed by \mathcal{SCND} placing n points, then \mathcal{FRST} has a winning strategy. In fact, if $Q = [0, 2n]$ and \mathcal{FRST} plays on the odd integer points $\{1, 3, 5, \dots, 2n - 1\}$, then \mathcal{SCND} 's payoff is less than $1/2$. So there is a set of “key points” in the one-dimensional arena, whose possession guarantees winning the game.

In this paper we show that no such set of “key points” exists in the two-dimensional case. More strongly, for each set \mathcal{F} of n points, there is a set \mathcal{S} of n \mathcal{SCND} points such that \mathcal{SCND} 's payoff is at least $1/2 + \alpha$, for an absolute constant $\alpha > 0$ and n large enough.

From now on, let Q be the square $[0, \sqrt{n}]^2$, of area n , so that the average area per \mathcal{FRST} 's point is 1. To win the game, \mathcal{SCND} needs to find n points such that their average area is at least $1/2 + \alpha$. We first show that it is very easy to find *one* such point—in fact, a *random* point in $Q \setminus \mathcal{F}$ has this property. Since this is the key idea of our proof, we first present it in a modified setting where the arena Q has the topology of a torus, eliminating boundary effects. We then proceed to prove this result for the square with its standard topology, showing how we can handle the square boundary, and proceed to prove the result for n \mathcal{SCND} points. Finally, we show that the result generalizes to higher dimensions as well.

2 The torus case

To present the (simple) main idea of our proofs in a setting free of technical complications due to effects near the boundary of Q , we assume in this section that the square Q has the topology of a torus. To be precise, we identify the left and right edges of Q , as well as the top and bottom edges, while retaining the normal Euclidean metric in Q .

Proposition 1 *There exist constants $\beta > 0$ and n_0 such that for every n -point set \mathcal{F} in the square arena Q with torus topology, $n \geq n_0$, there is a point $x \in Q \setminus \mathcal{F}$ with $\text{vol}(R(x, \mathcal{F})) \geq \frac{1}{2} + \beta$. In fact, x can be selected uniformly at random: $\mathbf{E}[\text{vol}(R(x, \mathcal{F}))] \geq \frac{1}{2} + \beta$, where $\mathbf{E}[\cdot]$ denotes expectation with respect to uniform random selection of $x \in Q$.*

Proof: If there is a point $p \in Q$ such that $\text{dist}(p, \mathcal{F}) > \sqrt{n}/4$, the proposition holds: With constant probability the point x will grab a constant fraction of n . If n is large enough,¹ this is more than, say, 1. In the following we can therefore assume that no such point p exists.

Let I_A denote the characteristic function of a set A .

$$\mathbf{E}[\text{vol}(R(x, \mathcal{F}))] = \frac{1}{\text{vol}(Q)} \int_Q \int_Q I_{R(x, \mathcal{F})}(y) \, dy \, dx$$

¹This is the only restriction on n in this proof. When we start to take boundary effects into account, we will have to assume n to be larger by several orders of magnitude.

$$= \frac{1}{n} \int_Q \text{vol}(\{x \in Q : y \in R(x, \mathcal{F})\}) dy$$

by Fubini's theorem.

A point $y \in Q$ lies in $R(x, \mathcal{F})$ if and only if $\text{dist}(y, x) \leq r = \text{dist}(y, \mathcal{F})$, and so the set $\{x \in Q : y \in R(x, \mathcal{F})\} = \{x \in Q : \text{dist}(x, y) \leq r\}$. Since $r \leq \sqrt{n}/4$, this is a disc of radius r centered at y (possibly wrapping around the edges of Q).

Our integral becomes $\frac{\pi}{n} \int_Q \text{dist}(y, \mathcal{F})^2 dy$, a quantity that we denote by $F_0(\mathcal{F})$. We split it into integrals over \mathcal{F} 's Voronoi cells: $F_0(\mathcal{F}) = \frac{\pi}{n} \sum_{w \in \mathcal{F}} \int_{\text{cell}_{\mathcal{F}}(w)} \text{dist}(y, w)^2 dy$, where $\text{cell}_{\mathcal{F}}(w)$ is the region of w in the Voronoi diagram of \mathcal{F} restricted to Q .

Among all convex bodies $C \subset \mathbf{R}^2$ of area a , the integral $\int_C \text{dist}(y, w)^2 dy$ is minimized by the disc C_0 of area a centered at w (somewhat informally, moving a piece of C closer to w decreases the integral, and such a move is possible for any C but that disc). Moreover, for later use we note that if C is a convex k -gon then $\int_C \text{dist}(y, w)^2 dy \geq (1 + \varepsilon_k) \int_{C_0} \text{dist}(y, w)^2 dy$ with a suitable small $\varepsilon_k > 0$.

The value of that integral over such a disc is

$$\int_{C_0} \text{dist}(y, w)^2 dy = \int_0^{\sqrt{a/\pi}} r^2 \cdot 2\pi r dr = \frac{a^2}{2\pi}.$$

Let's use the notation $a_w = \text{vol}(\text{cell}_{\mathcal{F}}(w))$. Then,

$$F_0(\mathcal{F}) = \frac{\pi}{n} \sum_{w \in \mathcal{F}} \int_{\text{cell}_{\mathcal{F}}(w)} \text{dist}(y, w)^2 dy \geq \frac{1}{2n} \sum_{w \in \mathcal{F}} a_w^2 \geq \frac{1}{2n} \frac{\left(\sum_{w \in \mathcal{F}} a_w\right)^2}{n} \geq \frac{1}{2}$$

by Cauchy-Schwarz.

So we see that for a random point x , the expected region size is at least $\frac{1}{2}$, but we want $\frac{1}{2} + \beta$. By the remark above, if $\text{cell}_{\mathcal{F}}(w)$ has at most k sides then $\int_{\text{cell}_{\mathcal{F}}(w)} \text{dist}(y, w)^2 dy \geq (1 + \varepsilon_k) \cdot \frac{a_w^2}{2\pi}$. Let $\mathcal{F}_f \subseteq \mathcal{F}$ consist of the points whose regions in the Voronoi diagram of \mathcal{F} have fewer than 12 sides. Since the average number of sides of a region in a planar Voronoi diagram is below 6 (using planarity of the dual graph, the Delaunay triangulation), we have $|\mathcal{F}_f| \geq \frac{1}{2}n$.

So we win the factor $1 + \varepsilon_{11}$ in at least half of the regions and lose nothing in the other regions. The only problem is that the regions of \mathcal{F}_f could together occupy only a tiny fraction of the area of Q and then this win would not reach the threshold $\beta > 0$ that we seek. But if they occupy, say, less than $\frac{1}{4}$ of the total area then the average area of the remaining regions (of $\mathcal{F} \setminus \mathcal{F}_f$) is at least $\frac{3}{2}$ (at most $\frac{1}{2}n$ regions take up area at least $\frac{3}{4}n$). Then the Cauchy-Schwarz inequality used in the calculation above becomes strict and we win a constant factor in the regions of $\mathcal{F} \setminus \mathcal{F}_f$. \square

3 The proof with boundary effects

The torus arena conveniently removed the need to consider the boundary effects. We now prove the same result for the square with boundary

Proposition 2 *There exist constants $\beta > 0$ and n_0 such that for every n -point set $W \subset Q$, $n \geq n_0$, we have $\mathbf{E}[\text{vol}(R(x, \mathcal{F}))] \geq \frac{1}{2} + \beta$.*

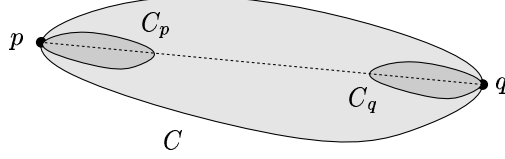


Figure 1: At least $1/16$ of the area of C is not covered by $B(w, \frac{1}{4}D)$.

Proof: As in the proof of Proposition 1, we can rewrite the expected area as

$$\begin{aligned}
F(\mathcal{F}) &= \frac{1}{n} \int_Q \text{vol}(\{x \in Q : y \in R(x, \mathcal{F})\}) \, dy \\
&= \frac{1}{n} \int_Q \text{vol}(B(y, \text{dist}(y, \mathcal{F})) \cap Q) \, dy \\
&= \frac{1}{n} \sum_{w \in \mathcal{F}} \int_{\text{cell}_{\mathcal{F}}(w)} \text{vol}(B(y, \text{dist}(y, w)) \cap Q) \, dy,
\end{aligned}$$

where $B(x, r)$ is the disc of radius r centered at x . We wish to bound $F(\mathcal{F})$ from below by $\frac{1}{2} + \beta$.

Let's choose a large constant D (the requirements on D will become apparent later). We call a region $\text{cell}_{\mathcal{F}}(w)$ *long* if it has diameter at least D and *short* otherwise, and we denote by \mathcal{F}_ℓ and \mathcal{F}_s the subsets of \mathcal{F} corresponding to the long and short regions, respectively.

First we consider the long regions. We note that for any $w, y \in Q$,

$$\text{vol}(B(y, \text{dist}(y, w)) \cap Q) \geq \frac{1}{2} \cdot \text{dist}(y, w)^2 \quad (1)$$

(the extreme case is w and y in opposite corners of Q).

Now let $w \in \mathcal{F}_\ell$ and write $C = \text{cell}_{\mathcal{F}}(w)$. We claim that at least $\frac{1}{16}$ of the area of C lies at distance at least $\frac{1}{4}D$ from w ; in other words, $\text{vol}(C \setminus B(w, \frac{1}{4}D)) \geq \frac{1}{16}a_w$ (the constant can be improved). Let p, q be a diametrical pair of points of C , and place two copies C_p, C_q of $C/4$ inside C so that they share a common tangent to C at p and q , respectively, where $C/4$ is the shape resulting from shrinking C by a factor of 4. Clearly, the distance between C_p and C_q is $D/2$, and consequently, either C_p or C_q do not intersect $B = B(w, \frac{1}{4}D)$. Thus, the area of C not covered by \mathcal{S} is at least $\text{vol}(C_p) = \text{vol}(C_q) = \text{vol}(C)/16$. See Fig. 1.

It follows that $\int_{\text{cell}_{\mathcal{F}}(w)} \text{vol}(B(y, \text{dist}(y, w)) \cap Q) \, dy \geq \frac{1}{2} \cdot \frac{D^2}{16} \cdot \frac{1}{16}a_w > \frac{D^2}{2000}a_w$ for every $w \in \mathcal{F}_\ell$, and so the contribution of the long regions to $F(\mathcal{F})$ is at least $\frac{D^2}{2000n}A_\ell$, where $A_\ell = \sum_{w \in \mathcal{F}_\ell} a_w$.

Next, we consider the short regions (of diameter $\leq D$), and among those only the *inner* ones, whose distance to the boundary of Q is at least D . Let \mathcal{F}_{si} be the corresponding subset of \mathcal{F} . We have $A_{si} = \sum_{w \in \mathcal{F}_{si}} a_w \geq n - 8D\sqrt{n} - A_\ell$. For the short inner regions, the disc $B(y, \text{dist}(y, w))$ lies completely inside Q and so their contribution to $F(\mathcal{F})$ behaves as in the proof of Proposition 1; it equals

$$\frac{\pi}{n} \sum_{w \in \mathcal{F}_{si}} \int_{\text{cell}_{\mathcal{F}}(w)} \text{dist}(y, w)^2 \, dy.$$

As we saw above, this quantity is bounded below by

$$\frac{1}{2n} \sum_{w \in \mathcal{F}_{si}} a_w^2 \geq \frac{1}{2n} \frac{A_{si}^2}{|\mathcal{F}_{si}|}.$$

Now we distinguish several cases depending on the orders of magnitude of A_ℓ and $|\mathcal{F}_{si}|$. First suppose that $A_\ell \geq \frac{n}{2D}$; then the contribution of A_ℓ alone suffices: $F(\mathcal{F}) \geq \frac{D^2}{2000n}A_\ell \geq \frac{D}{4000} > \frac{1}{2} + \beta$ for D large enough. Next, let $A_\ell < \frac{n}{2D}$, which for large n implies $A_{si} \geq (1 - \frac{1}{D})n$. Now two cases are distinguished according to $|\mathcal{F}_{si}|$. For $|\mathcal{F}_{si}| \leq (1 - \frac{4}{D})n$, we obtain

$$F(\mathcal{F}) \geq \frac{1}{2n} \frac{A_{si}^2}{|\mathcal{F}_{si}|} \geq \frac{(1 - \frac{1}{D})^2}{2(1 - \frac{4}{D})} \geq \frac{1}{2} + \frac{1}{D}$$

which is the desired bound.

Finally it remains to deal with the case $A_{si} \geq (1 - \frac{1}{D})n$ and $|\mathcal{F}_{si}| \geq (1 - \frac{4}{D})n$. If D is very large, we are essentially in the situation analyzed in the proof of Proposition 1 and practically the same argument shows that $F(\mathcal{F}) \geq \frac{1}{2} + \beta$ in this case as well (using the fact that $\frac{1}{D}$ is much smaller than ε_{11}). \square

4 The main result

A key ingredient in the proof of our main theorem is the following lemma, showing that if $SC\mathcal{N}\mathcal{D}$ throws in δn points at random, instead of one as in Proposition 2, then his expected area gain still exceeds $\frac{1}{2}\delta n$ at least by a fixed fraction, provided that $\delta > 0$ is sufficiently small.

Lemma 3 *There exist constants $\beta_1 > 0$, $\delta > 0$, and n_0 such that for every n -point set $\mathcal{F} \subset Q$, $n \geq n_0$, if $S \subset Q$ is obtained by δn independent random draws from the uniform distribution on Q then $\mathbf{E}[\text{vol}(R(S, \mathcal{F}))] \geq (\frac{1}{2} + \beta_1)\delta n$.*

Moreover, for every given sufficiently large constant D , the constants δ and n_0 above can be adjusted so that whenever the total area A_ℓ of the long regions (of diameter at least D) exceeds $\frac{n}{2D}$ then $\mathbf{E}[\text{vol}(R(S, \mathcal{F}))] \geq 2\delta n$.

Proof:(Sketch) This is very similar to the proof of Proposition 2. Intuitively, for small δ , the δn independent random points are likely to interact very little and their expected area gain is likely to be nearly $(\delta - O(\delta^2))n$ times the expected area gain of a single point.

This time we have

$$\mathbf{E}[\text{vol}(R(S, \mathcal{F}))] = \int_Q \text{Prob}[y \in R(S, \mathcal{F})] dy.$$

Here $P(y) = \text{Prob}[y \in R(S, \mathcal{F})]$ is the probability with respect to the random choice of the set S . Namely,

$$\begin{aligned} P(y) &= \text{Prob}[S \cap B(y, \text{dist}(y, \mathcal{F})) \neq \emptyset] \\ &= 1 - (\text{Prob}[x \notin B(y, \text{dist}(y, \mathcal{F}))])^{\delta n} \\ &= 1 - (1 - \frac{1}{n} \cdot \text{vol}(B(y, \text{dist}(y, \mathcal{F})) \cap Q))^{\delta n}. \end{aligned}$$

Let's write $\rho(y) = \frac{1}{n} \cdot \text{vol}(B(y, \text{dist}(y, \mathcal{F})) \cap Q)$. If y lies in a short region of the Voronoi diagram of \mathcal{F} then $\rho(y) \leq \frac{C_D}{n}$ with C_D depending on D only, and δC_D can be made as small as desired by choosing δ sufficiently small. Then we obtain $P(y) = 1 - (1 - \rho(y))^{\delta n} \geq \delta n \rho(y) + O((\delta n \rho(y))^2) \geq$

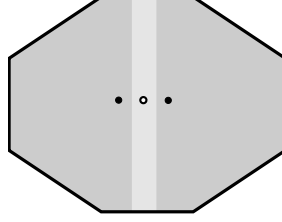


Figure 2: Two points of $SC\mathcal{N}\mathcal{D}$ can take over almost a complete cell of $\mathcal{F}\mathcal{R}\mathcal{S}\mathcal{T}$.

$\delta n \rho(y) \cdot (1 - \gamma)$ with γ a small constant. Thus, the contribution of a short Voronoi region to $\mathbf{E}[\text{vol}(R(\mathcal{S}, \mathcal{F}))]$ is at least $(1 - \gamma)\delta n$ times the contribution of that region to the expected area gained by a single random point as in Proposition 2. All the calculations involving short regions can be done in exactly the same way. It remains to show that if the total area A_ℓ of the long regions is at least $\frac{n}{2D}$ then these regions contribute at least $2\delta n$ to $\mathbf{E}[\text{vol}(R(\mathcal{S}, \mathcal{F}))]$.

In the proof of Proposition 2, Eq. (1), we have shown $\rho(y) \geq \frac{1}{2n} \cdot \text{dist}(y, w)^2$ for $y \in \text{cell}_{\mathcal{F}}(w)$. We also know that $\text{dist}(y, w) \geq \frac{1}{4}D$ for y in at least $\frac{1}{16}$ of the area of each long region. For these y , we have $P(y) \geq 1 - e^{-\rho(y)\delta n} \geq 1 - e^{-D^2\delta/200} \geq \frac{D^2\delta}{400}$ (assuming $\delta < D^{-2}$). The whole integral over all the long regions is then at least $\frac{1}{16}A_\ell \geq \frac{n}{32D}$ times this quantity and therefore larger than $2\delta n$ with ample room to spare. \square

We can now prove our main theorem.

Theorem 4 *There exist constants $\alpha > 0$ and n_0 such that for every $n \geq n_0$, $SC\mathcal{N}\mathcal{D}$ can always win at least $\frac{1}{2} + \alpha$ in the Voronoi game. That is, for every n -point set $\mathcal{F} \subset Q$ there exists an n -point set $\mathcal{S} \subset Q \setminus \mathcal{F}$ with $\text{vol}(R(\mathcal{S}, \mathcal{F})) \geq (\frac{1}{2} + \alpha)\text{vol}(Q)$.*

Proof: Let $w \in \mathcal{F}$. A takeover of w 's region means that $SC\mathcal{N}\mathcal{D}$ places two of his points very close to w with w as the center of symmetry. See Fig. 2. In this way, he captures almost all of $\text{cell}_{\mathcal{F}}(w)$. This suggests the following strategy for $SC\mathcal{N}\mathcal{D}$: A takeover of the $\frac{1}{2}n$ largest $\mathcal{F}\mathcal{R}\mathcal{S}\mathcal{T}$ regions guarantees $SC\mathcal{N}\mathcal{D}$ a payoff arbitrarily close to $\frac{1}{2}n$. This does not prove the theorem, in general, but it fails to do so only if almost all of $\mathcal{F}\mathcal{R}\mathcal{S}\mathcal{T}$'s regions have almost the same area. Thus, if more than εn $\mathcal{F}\mathcal{R}\mathcal{S}\mathcal{T}$ regions have area below $1 - \varepsilon$, for some constant $\varepsilon > 0$, then the takeover strategy implies the theorem. It therefore suffices to describe a strategy for $SC\mathcal{N}\mathcal{D}$ when all but εn of $\mathcal{F}\mathcal{R}\mathcal{S}\mathcal{T}$'s regions have area at least $1 - \varepsilon$. (A similar trick would also simplify the proof of Proposition 2 if we didn't want to prove the claim about a random point but only the existence of a point capturing at least $\frac{1}{2} + \beta$.)

First $SC\mathcal{N}\mathcal{D}$ chooses a set \mathcal{S}_0 of δn points as in Lemma 3; that is, with $\text{vol}(R(\mathcal{S}_0, \mathcal{F})) \geq (1 + \beta_1)\delta n$ and even with $\text{vol}(R(\mathcal{S}_0, \mathcal{F})) \geq 2\delta n$ if $A_\ell \geq \frac{n}{2D}$.

If $A_\ell \geq \frac{n}{2D}$ then $\mathcal{F}\mathcal{R}\mathcal{S}\mathcal{T}$ now has n regions of total area $A_{\mathcal{F}} \leq (1 - 2\delta)n$ and $SC\mathcal{N}\mathcal{D}$ still has $(1 - \delta)n$ points to play. He takes over the $\frac{1}{2}(1 - \delta)n$ largest among the current regions of $\mathcal{F}\mathcal{R}\mathcal{S}\mathcal{T}$. In this way, $SC\mathcal{N}\mathcal{D}$ has captured at least area arbitrarily close to

$$n - A_{\mathcal{F}} + \frac{1}{2}(1 - \delta)n \cdot \frac{A_{\mathcal{F}}}{n} = n - \frac{1}{2}(1 + \delta)A_{\mathcal{F}} > \frac{1}{2}(1 + \delta)n.$$

Next, we suppose that $A_\ell < \frac{n}{2D}$. Let's consider a point $w \in \mathcal{F}_s$ defining a short region and call w contaminated if $SC\mathcal{N}\mathcal{D}$ has captured some point of $\text{cell}_{\mathcal{F}}(w)$ by the set \mathcal{S}_0 . A short region can

only be contaminated by a point $b \in \mathcal{S}_0$ if $\text{dist}(b, w) \leq 2D$. Therefore, the total area of contaminated short regions is $O(D^2 \delta n) < \frac{n}{3}$, say, and so regions of total area at least $\frac{n}{2}$ remain uncontaminated. Now we use the assumption that all but εn of \mathcal{FRT} 's regions have area at least $1 - \varepsilon$. \mathcal{SCND} can now overtake the $\frac{1}{2}(1 - \delta)n$ largest uncontaminated regions. This implies that the number of uncontaminated regions of size $\geq 1 - \varepsilon$ is at least $n/2 - \varepsilon n$. Thus, \mathcal{SCND} can now occupy at least $\min(n/2 - \varepsilon n, (1 - \delta)n/2) \geq \frac{1}{2}(1 - \delta)n - \varepsilon n$ cells, to gain total area at least

$$\left(\frac{1}{2} + \beta_1\right)\delta n + \left(\frac{1}{2}(1 - \delta)n - \varepsilon n\right)(1 - \varepsilon) = \left(\frac{1}{2} + \beta_1\right)\delta n + \frac{1}{2}(1 - \delta - 2\varepsilon)(1 - \varepsilon)n.$$

If ε is very small compared to δ and β_1 then this is at least $(\frac{1}{2} + \alpha)n$ with α close to $\beta_1\delta$. This concludes the proof of the theorem. \square

5 The higher-dimensional case

The proof of Proposition 1 (and therefore of Lemma 3) exploited the fact that the Voronoi diagram is a planar graph, and therefore at least half of all Voronoi cells have at most 11 edges. This is not true in dimensions higher than two. We used this fact to argue that Voronoi cells cannot be arbitrarily similar to a disc. To extend our theorem to higher dimensions, we need to give a more general proof of this claim.

Definition 5 A convex body C with a point p is a μ -ball, for $\mu > 0$, if there exists a radius $r > 0$ such that $B(p, r) \subseteq C \subseteq B(p, r(1 + \mu))$.

Lemma 6 If C is a convex body in \mathbb{R}^d with center p which is not a μ -ball, then

$$\int_C c_d \text{dist}(y, p)^d dy \geq (1 + \beta)L \quad \text{where} \quad L = \int_D c_d \text{dist}(y, p)^d dy = \frac{\text{vol}(C)^2}{2},$$

where D is a ball of the same volume as C centered at p , $\beta > 0$ is a constant that depends only on μ and d , and c_d is the volume of the unit ball in \mathbb{R}^d .

Proof: Let $D = B(p, R)$, where $R = (\text{vol}(C)/c_d)^{1/d}$. Then,

$$L = \int_D c_d \text{dist}(y, p)^d dy = \int_{r=0}^R \left((c_d r^d) \cdot (c_d dr^{d-1}) \right) dr = \frac{c_d^2}{2} R^{2d} = \frac{\text{vol}(C)^2}{2}.$$

As for the other claim, let r', R' be the largest (resp. smallest) radius so that $B(p, r') \subseteq C \subseteq B(p, R')$. Since C is not a μ -ball, it follows that there exists a positive constant β_1 such that $R(1 + \beta_1) \leq R'$. In particular, this implies that there exists a constant β_2 , such that $\text{vol}(K) \geq \beta_2 \text{vol}(C)$, where $K = C \setminus B(p, R(1 + \beta_1/4))$. Namely,

$$\begin{aligned} \int_C c_d \text{dist}(y, p)^d dy &\geq \int_D c_d \cdot \text{dist}(y, p)^d dy + \int_K c_d \left(\left(R \cdot (1 + \beta_1/4) \right)^d - R^d \right) dy \\ &\geq (1 + \beta) \int_D c_d \cdot \text{dist}(y, p)^d dy = (1 + \beta)L, \end{aligned}$$

where $\beta > 0$ is an appropriate constant that depends only on d and μ . \square

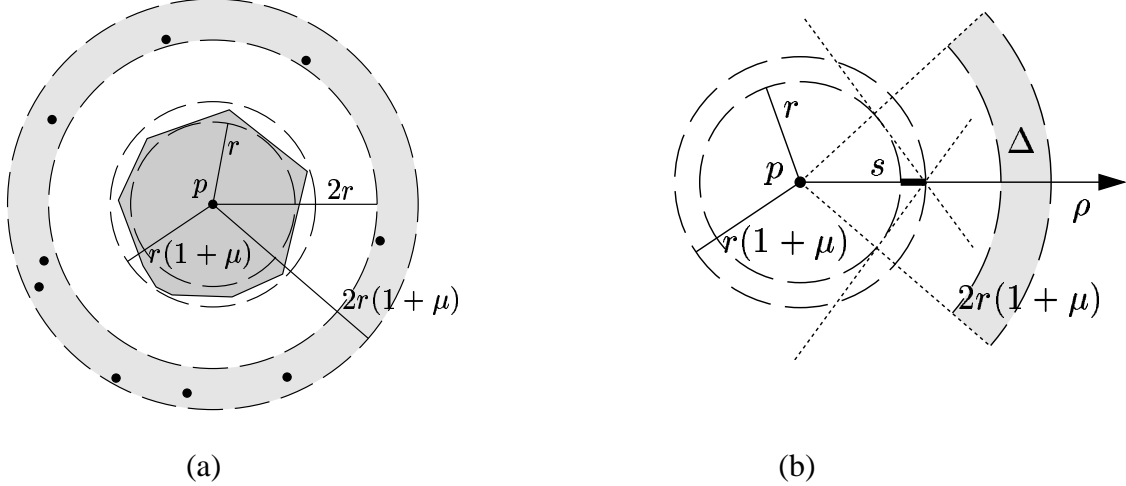


Figure 3: (a) For a μ -ball Voronoi cell C_p all neighboring sites must be in a spherical shell around it. (b) The sites of $N(p)$ are densely spread in this spherical shell as there must be a site inside the intersection Δ between the spherical shell and a cone of angular radius $4\sqrt{\mu}$.

Lemma 7 *Let Q be a hypercube in \mathbb{R}^d , and let P be a set of points in Q . Let $V(P)$ denote the decomposition of Q into convex cells by the Voronoi diagram of P restricted to Q . Then there exists a constant $\mu > 0$, which depends only on d , so that the total volume of the μ -ball cells in $V(P)$ is bounded by $\text{vol}(Q)/2$.*

Proof: Consider a cell C_p of $V(P)$ that is a μ -ball, and let $p \in P$ be its center. Let $B(p, r)$ be the largest ball that is contained inside C_p . Let $N(p)$ be the set of points of P whose Voronoi cells have a common boundary with C_p .

Observe that the distance of any point of $N(p)$ to p is at least $2r$ and at most $2r(1 + \mu)$. Furthermore, any angular cone of angular angle $4\sqrt{\mu}$ emanating from p must include a point of $N(p)$. Indeed, consider such a cone Z with a ray ρ as its rotational axis and angular radius $4\sqrt{\mu}$, where ρ emanates from p . Let s denote the intersection of ρ with the spherical shell $B(p, r(1 + \mu)) \setminus B(p, r)$. Since one endpoint of s is outside C_p , and the other is inside C_p , it follows that there must be a point $q \in P$, so that the bisector of p and q intersects s . It is now straightforward to verify that q is inside Z . See Figure 3.

This implies that $N(p)$ is dense. Indeed, consider a point $q \in N(p)$. Its nearest point in $N(p)$ is at distance at most $2r(1 + \mu) \cdot 2 \cdot 4\sqrt{\mu} = O(r\sqrt{\mu})$. On the other hand, the Voronoi cell C_q of q has a point on its boundary of distance $\geq r$ from q (as it shares a boundary point u with C_p , $\text{dist}(u, p) \geq r$, and $\text{dist}(u, q) = \text{dist}(u, p)$). (This also implies that C_p is not adjacent to the boundary of Q .)

That is, C_q is not a γ -ball, where $\gamma = \Omega(r/r\sqrt{\mu} - 1) = \Omega(1/\sqrt{\mu})$. Thus, by making μ small enough, we can ensure that C_q is not a μ -ball.

We have shown that every μ -ball cell in $V(P)$ is surrounded by cells that are not μ -balls. We will charge the volume of such a μ -ball to its surrounding cells as follows. For a point $p \in P$ whose Voronoi cell is a μ -ball, let r_p be the radius of the largest ball contained inside C_p centered at p , and let $U_p = B(p, 1.8r_p)$ be the *region of influence* of p . Clearly, $U_p \cap P = \{p\}$, $\text{vol}(U_p) \geq (1.8/(1 + \mu))^d \text{vol}(C_p) \geq 2\text{vol}(C_p)$, for μ small enough. By picking μ small enough, we can also guarantee that the regions of influence of the μ -balls of $V(P)$ are disjoint. We charge the volume

of a μ -ball to its region of influence, establishing the claim. \square

Plugging Lemmas 6 and 7 into the proof of Theorem 4 gives us the following result. The details of the adaption are straightforward and omitted.

Theorem 8 *There exist constants $\alpha > 0$ and n_0 depending only on the dimension d , such that for every $n \geq n_0$, $SC\mathcal{N}\mathcal{D}$ can always win at least $\frac{1}{2} + \alpha$ in the Voronoi game with a hypercube playing arena Q in \mathbb{R}^d . That is, for every n -point set $\mathcal{F} \subset Q$ there exists an n -point set $S \subset Q \setminus \mathcal{F}$ with $\text{vol}(R(S, \mathcal{F})) \geq (\frac{1}{2} + \alpha)\text{vol}(Q)$.*

6 Conclusions and open problems

We considered the Voronoi game on a square or hypercube board Q , played in a single round: \mathcal{FRST} starts by placing n points \mathcal{F} in Q , then $SC\mathcal{N}\mathcal{D}$ places another n points S disjoint from \mathcal{F} , and finally the winner is determined.

Our considerations appear to generalize without much change to sufficiently “fat” convex playing arenas in the plane. On the other hand, when the playing arena degenerates to a line segment, we have reached the one-dimensional case where \mathcal{FRST} , not $SC\mathcal{N}\mathcal{D}$, has a winning strategy [ACC⁺01]. It would be interesting to understand the behavior of the game with a rectangular playing arena as a function of the aspect ratio of the rectangle.

What happens when the number of points played by \mathcal{FRST} and $SC\mathcal{N}\mathcal{D}$ are not identical? Specifically, let λ be a real number between 0 and 2. Consider the game where \mathcal{FRST} plays n points and $SC\mathcal{N}\mathcal{D}$ plays λn points. Let $f(\lambda, n)$ be the payoff to $SC\mathcal{N}\mathcal{D}$ in this Voronoi game. It is not hard to show that $f(0, n) = 0$ and that $\lim_{n \rightarrow \infty} f(2, n) = 1$. We know that $f(\lambda, n) > (\frac{1}{2} + \epsilon)\lambda$ for some positive ϵ and n large enough, as long as λ is bounded away from 0 and 2. It would be interesting to get a better idea of the behavior of f . Does $\lim_{n \rightarrow \infty} f(\lambda, n)$ exist for all λ ?

We have shown that for any set of n \mathcal{FRST} points, there is a $SC\mathcal{N}\mathcal{D}$ point that grabs a “large” Voronoi cell. It would be interesting to find configurations of the \mathcal{FRST} points for which no $SC\mathcal{N}\mathcal{D}$ point can do too well. Obvious candidates are grid arrangements of the \mathcal{FRST} points, such as the square grid or hexagonal grid.

The original version of the Voronoi game [ACC⁺01] is played in more than one round: \mathcal{FRST} and $SC\mathcal{N}\mathcal{D}$ alternate placing points on the board Q . The value of this game and the optimal strategies are still unknown for dimension higher than one. It is easy to verify that if the given board Q is symmetric, but has no center of symmetry, then $SC\mathcal{N}\mathcal{D}$'s payoff is at least $\frac{1}{2}$. This can be guaranteed by responding to each move of \mathcal{FRST} with a point placed in the symmetric location. Many obvious questions remain open: Can $SC\mathcal{N}\mathcal{D}$ actually win the game for large n ? What happens with asymmetric boards?

References

- [ACC⁺01] H. Ahn, S. Cheng, O. Cheong, M. Golin, and R. van Oostrum. Competitive facility location along a highway. In *7th Annual International Computing and Combinatorics Conference*, volume 2108 of *LNCS*, pages 237–246, 2001.

- [Aur91] F. Aurenhammer. Voronoi diagrams: A survey of a fundamental geometric data structure. *ACM Comput. Surv.*, 23:345–405, 1991.
- [EL89] H.A. Eiselt and G. Laporte. Competitive spatial models. *European Journal of Operational Research*, 39:231–242, 1989.
- [ELT93] H.A. Eiselt, G. Laporte, and J.-F. Thisse. Competitive location models: A framework and bibliography. *Transportation Science*, 27:44–54, 1993.
- [TFM89] R.L. Tobin, T.L. Friesz, and T. Miller. Existence theory for spatially competitive network facility location models. *Annals of Operations Research*, 18:267–276, 1989.